

# Growing Hypervisor 0day with Hyperseed

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# whoami

- Daniel King
  - MSRC V&M
  - Keen Security Lab
  - Member of "Master of Pwn" 2016
  - MSRC Nano Server Bounty Winner
  - Ruxcon 2016
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  - USAF
  - Matasano
  - RPISEC

# Overview

- Motivation
- Hyper-V Architecture
- Hypercalls
- Hyperseed
- Mutations
- Findings

# Motivation – Why Hypercalls?

## Importance

- Cornerstone of Hyper-V

## Accessibility

- Accessible from kernel (Ring-0) mode

## Simplicity

- Well-documented, good starting point

## Complexity

- Diverse input & output format

# Motivation – Bug Bounty

## RCE w/ Exploit (Guest-to-Host-Escape)

- \$250,000 (Hypervisor/Kernel) / \$150,000 (User-mode)

## RCE (Guest-to-Host-Escape)

- \$200,000 (Hypervisor/Kernel) / \$100,000 (User-mode)

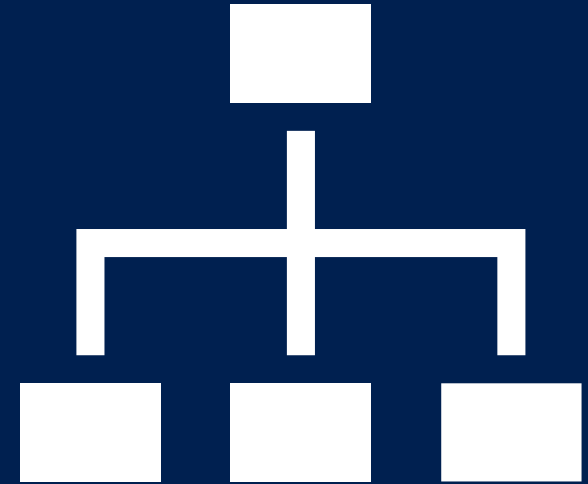
## Information Disclosure

- \$25,000 (Hypervisor/Kernel), \$15,000 (User-mode)

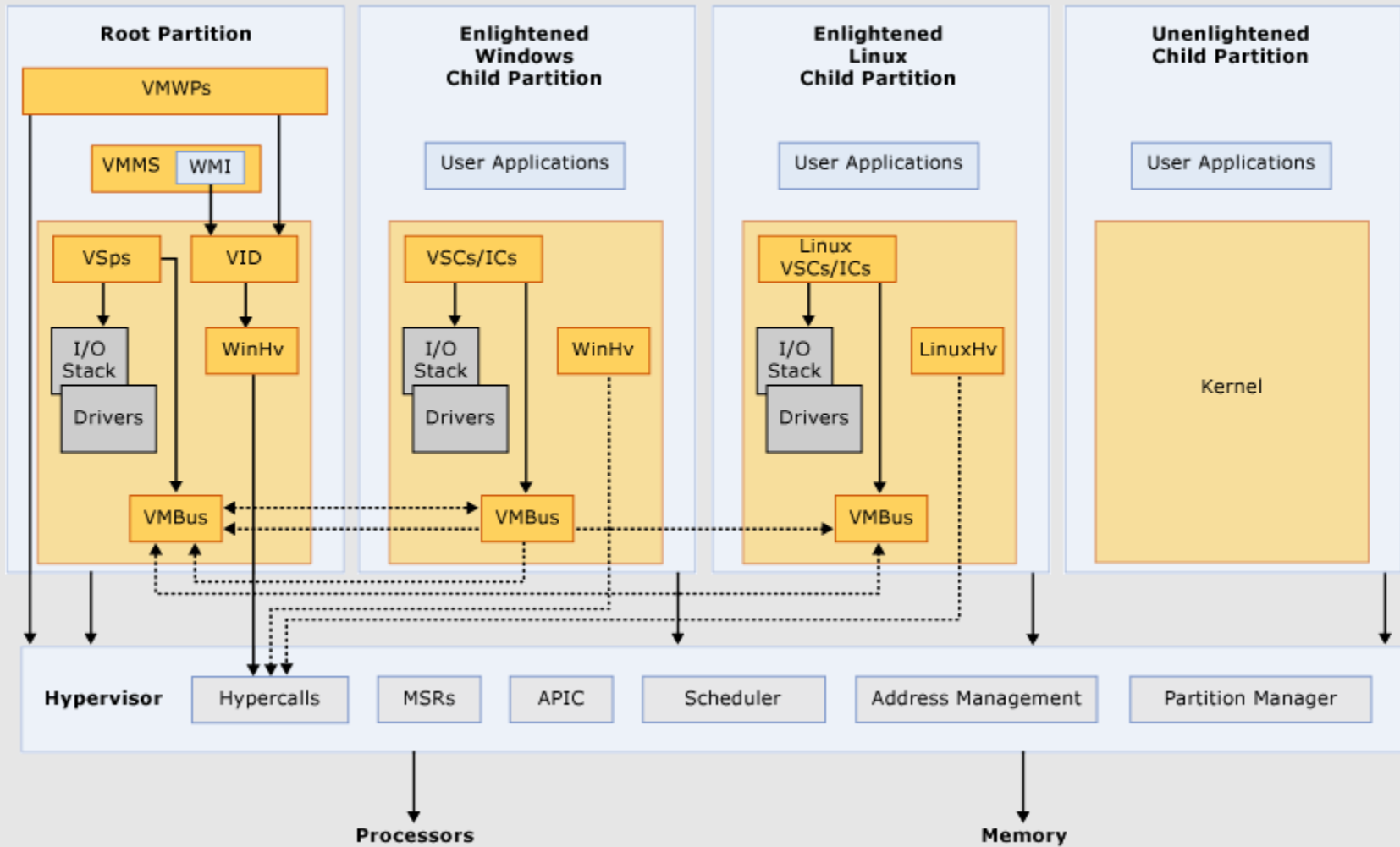
## Denial of Service

- \$15,000 (Hypervisor/Kernel)

# Hyper-V Architecture



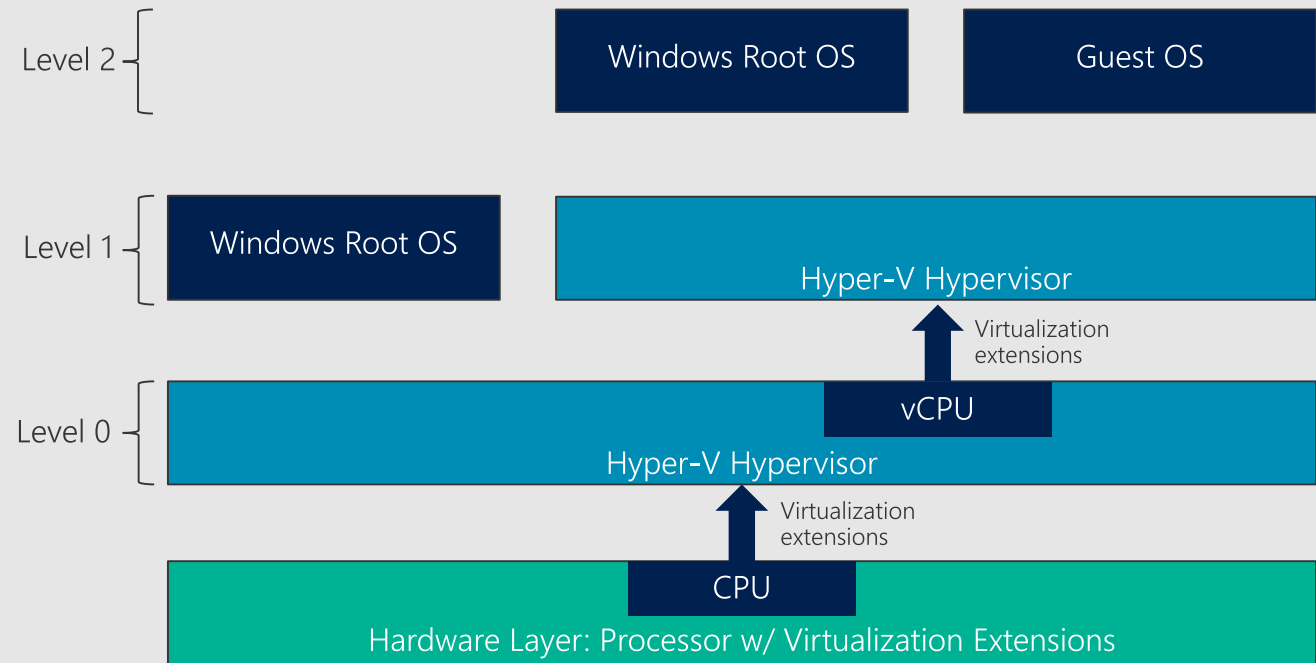
## Hyper-V High Level Architecture



<https://docs.microsoft.com/en-us/virtualization/hyper-v-on-windows/reference/hyper-v-architecture>

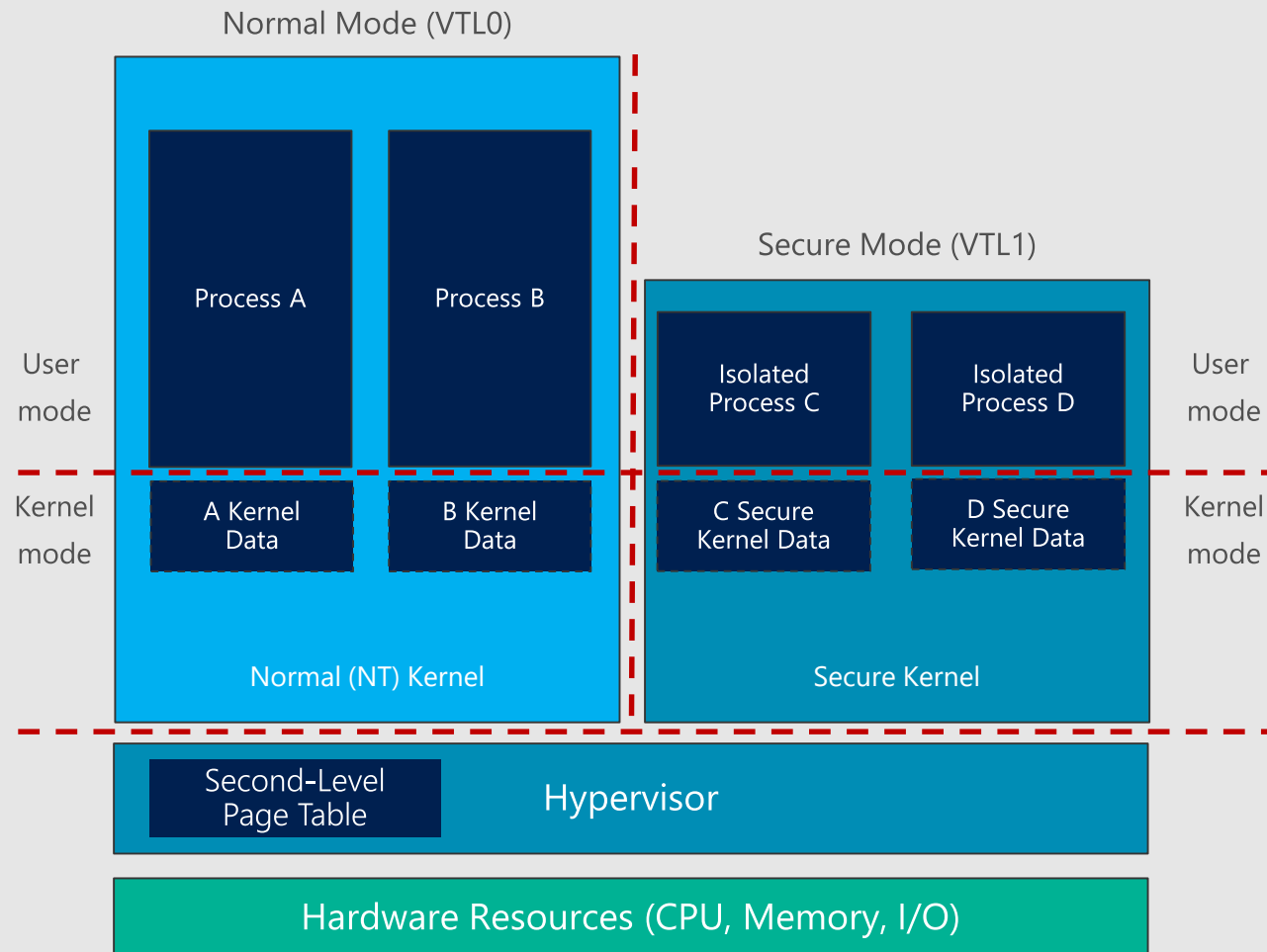
# Nested Virtualization

- L0 Hyper-V Hypervisor
- L1 Root
- L1 Guest
- L1 Hyper-V Hypervisor
- L2 Root
- L2 Guest





# Virtualization Based Security



- VBS introduces concept of trust levels
  - VTL0 – normal mode
  - VTL1 – secure (privileged) mode
- VTL1 controls which resources VTL0 can access
  - Memory, privileged CPU registers
- Hypervisor leverages hardware virtualization to enforce access permissions
  - SLAT – memory access protection for CPU
  - IOMMU – memory access protection for devices
- VTL1 is NOT a separate VM
  - VTL1 and VTL0 share memory, CPU's
  - Secure kernel coordinates with normal kernel for memory management, scheduling, etc.



# Hypercalls



# Hypercall Basics

- Calling mechanism for guests to interface with the hypervisor
  - Three calling conventions
    - Slow
    - Fast
    - XMM Fast
- Hypercall classes
  - Simple
  - Rep "repeat"
- Variable Sized Hypercalls
- Extended Hypercall Interface
  - Call codes above 0x8000

```
typedef union _HV_HYPERCALL_INPUT
{
    struct
    {
        UINT32 CallCode           : 16;
        UINT32 IsFast             : 1;
        UINT32 VariableHeaderSize : 9;
        UINT32 Reserved1          : 6;
        UINT32 CountOfElements    : 12;
        UINT32 Reserved2          : 4;
        UINT32 RepStartIndex       : 12;
        UINT32 Reserved3          : 4;
    };
    UINT64 AsUINT64;
} HV_HYPERCALL_INPUT, *PHV_HYPERCALL_INPUT;
```

```
typedef union _HV_HYPERCALL_OUTPUT
{
    struct
    {
        UINT16 CallStatus;
        UINT16 Reserved1;
        UINT32 ElementsProcessed : 12;
        UINT32 Reserved2        : 20;
    };
    UINT64 AsUINT64;
} HV_HYPERCALL_OUTPUT, *PHV_HYPERCALL_OUTPUT;
```

# How To Issue a Hypercall

- Hypercalls must be invoked from the most privileged guest processor mode, CPL0 / Kernel mode
- All hypercalls should be invoked through the architecturally-defined hypercall interface
  - vmcall instruction for Intel VMX
  - vmxcall instruction for AMD SVM
- Hypercall Page: abstract the difference between different virtualization implementations
  - `#define HV_X64_MSR_HYPERCALL 0x40000001`

# Establishing the Hypercall Interface

- Write guest id to HV\_X64\_MSR\_GUEST\_OS\_ID
- Allocate a page of memory for the hypercall page
- Write to the hypercall MSR HV\_X64\_MSR\_HYPERCALL
  - Set the Enable Hypercall Page bit
  - Write the GPA of the hypercall page

```
/*  
 * Setup the hypercall page and enable hypercalls.  
 * 1. Register the guest ID  
 * 2. Enable the hypercall and register the hypercall page  
 */  
guest_id = generate_guest_id(0, LINUX_VERSION_CODE, 0);  
wrmsrl(HV_X64_MSR_GUEST_OS_ID, guest_id);  
  
hypercall_pg = __vmalloc(PAGE_SIZE, GFP_KERNEL, PAGE_KERNEL_RX);  
  
if (hypercall_pg == NULL) {  
    wrmsrl(HV_X64_MSR_GUEST_OS_ID, 0);  
    return;  
}  
  
rdmsrl(HV_X64_MSR_HYPERCALL, hypercall_msr.as_uint64);  
hypercall_msr.enable = 1;  
hypercall_msr.guest_physical_address = vmalloc_to_pfn(hypercall_pg);  
wrmsrl(HV_X64_MSR_HYPERCALL, hypercall_msr.as_uint64);
```

FROM: <https://github.com/LIS/lis-next>

# Hypercall Page

**HYPERCALL:**

```
vmcall  
ret
```

Calling here will issue a hypercall

**VTL\_CALL:**

```
mov rax, rcx  
mov rcx, 0x0011  
vmcall  
ret
```

Calling here will issue a secure call from VTL0 to VTL1

**VTL\_RETURN:**

```
mov rax, rcx  
mov rcx, 0x0012  
vmcall  
ret
```

Calling here will issue a normal call from VTL1 to VTL0

```
nop  
nop  
.....  
nop
```

# HvCallVtlCall & HvCallVtlReturn

VTL 0 -> VTL 1

- nt!VslpEnterlumSecureMode
- vmcall(0x11)

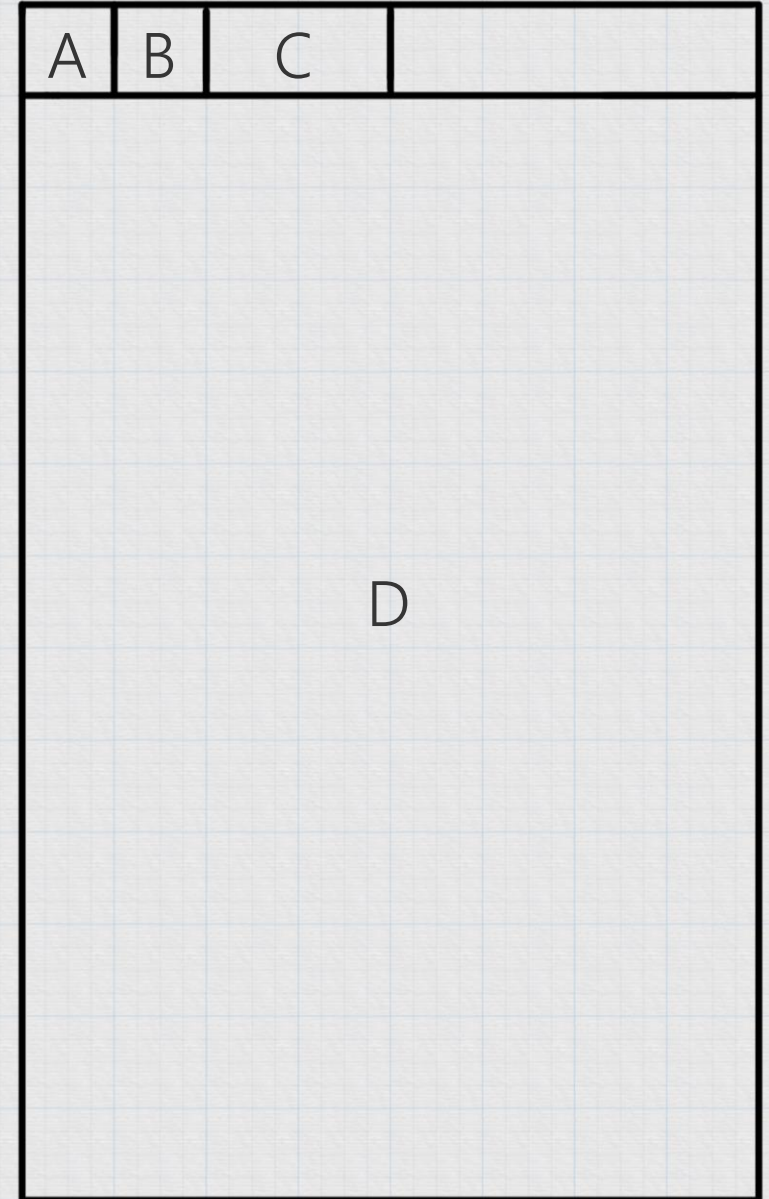


VTL 1 -> VTL 0

- sk!SkCallNormalMode
- vmcall(0x12)

# VTL Call & Return Buffer

- A [uint8\_t]: SecureOperation
  - 2: Invoke Secure Service
- B [uint8\_t]: NormalOperation
  - 0: Invoke Normal Service
- C [uint16\_t]: ServiceCode
- D [uint64\_t]: Parameters[12]





# Restrictions Enforced by Hypercall Handlers

# Restrictions Enforced by Hypercall Handlers

- The calling partition must possess a particular privilege

# Partition Properties

- Partitions are identified by using a partition ID
  - `#define HV_PARTITION_ID_INVALID ((HV_PARTITION_ID) 0x0)`
  - `#define HV_PARTITION_ID_SELF ((HV_PARTITION_ID) -1)`
- Each partition has a set of properties
  - Read a property of a partition by invoking `HvGetPartitionProperty`
  - Modify partitions property by invoking `HvSetPartitionProperty`
  - `HvPartitionPropertyPrivilegeFlags` defines what the partition is allowed to access
- Partition privilege flags represented by a bit-field

# Partition Privilege Flags

```
// Access to virtual MSRs
AccessVpRunTimeReg:1;
AccessPartitionReferenceCounter:1;
AccessSynicRegs:1;
AccessSyntheticTimerRegs:1;
AccessIntrCtrlRegs:1;
AccessHypercallMsrs:1;
AccessVpIndex:1;
AccessResetReg:1;
AccessStatsReg:1;
AccessPartitionReferenceTsc:1;
AccessGuestIdleReg:1;
AccessFrequencyRegs:1;
AccessDebugRegs:1;
AccessReenlightenmentControls:1;
Reserved1:18;
```

```
// Access to hypercalls
CreatePartitions:1;
AccessPartitionId:1;
AccessMemoryPool:1;
AdjustMessageBuffers:1;
PostMessages:1;
SignalEvents:1;
CreatePort:1;
ConnectPort:1;
AccessStats:1;
Reserved2:2;
Debugging:1;
CpuManagement:1;
Reserved:1;
Reserved:1;
Reserved:1;
AccessVsm:1;
AccessVpRegisters:1;
Reserved:1;
Reserved:1;
EnableExtendedHypercalls:1;
StartVirtualProcessor:1;
Reserved3:10;
```

# Partition Privilege Flags - Child

```
// Access to virtual MSRs
AccessVpRunTimeReg:1;
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AccessStatsReg:1;
AccessPartitionReferenceTsc:1;
AccessGuestIdleReg:1;
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AccessDebugRegs:1;
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Reserved:1;
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AccessVpRegisters:1;
Reserved:1;
Reserved:1;
EnableExtendedHypercalls:1;
StartVirtualProcessor:1;
Reserved3:10;
```

# Partition Privilege Flags - Root

```
// Access to virtual MSRs
AccessVpRunTimeReg:1;
AccessPartitionReferenceCounter:1;
AccessSynicRegs:1;
AccessSyntheticTimerRegs:1;
AccessIntrCtrlRegs:1;
AccessHypercallMsrs:1;
AccessVpIndex:1;
AccessResetReg:1;
AccessStatsReg:1;
AccessPartitionReferenceTsc:1;
AccessGuestIdleReg:1;
AccessFrequencyRegs:1;
AccessDebugRegs:1;
AccessReenlightenmentControls:1;
Reserved1:18;
```

```
// Access to hypercalls
CreatePartitions:1;
AccessPartitionId:1;
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Reserved:1;
Reserved:1;
EnableExtendedHypercalls:1;
StartVirtualProcessor:1;
Reserved3:10;
```

# Partition Privilege Check

```
//  
// Check for CpuManagement permission.  
//  
status = HV_STATUS_SUCCESS;  
privileges.AsUINT64 = 0;  
privileges.CpuManagement = TRUE;  
HT_CHK(status = VmCheckCurrentPartitionPrivileges(privileges));
```



# Restrictions Enforced by Hypercall Handlers

- The calling partition must possess a particular privilege
- The partition being acted upon must be in a particular state
- The partition must be either a parent or child

# Partition State Check

```
//  
// This function locates a partition with the specified partition Id and  
// references it only if its state matches one of the specified values.  
//  
HT_CHK(status = ObReferencePartitionById(InputHeader->PartitionId,  
                                           ObPartitionReferenceIfActive,  
                                           ObPartitionRightParent,  
                                           ObPartitionSecurityNormal,  
                                           &partition,  
                                           &partitionReference));
```

# Restrictions Enforced by Hypercall Handlers

- The calling partition must possess a particular privilege
- The partition being acted upon must be in a particular state
- The partition must be either a parent or child
- The partition must be the root

# Root Partition Check

```
//  
// This hypercall can only be made by the highest enabled VTL in the root partition.  
//  
partition = VmGetCurrentPartition();  
if (!partition->Settings.IsRootPartition)  
{  
    HT_ERR(status = HV_STATUS_ACCESS_DENIED);  
}  
HT_CHK(status = VmCheckCurrentVtlAuthority(VmRootPartition, VmGetPtHighestEnabledVtl()));
```

# Restrictions Enforced by Hypercall Handlers

- The calling partition must possess a particular privilege
- The partition being acted upon must be in a particular state
- The partition must be the root
- The partition must be either a parent or child
- The virtual processor must be in a particular state

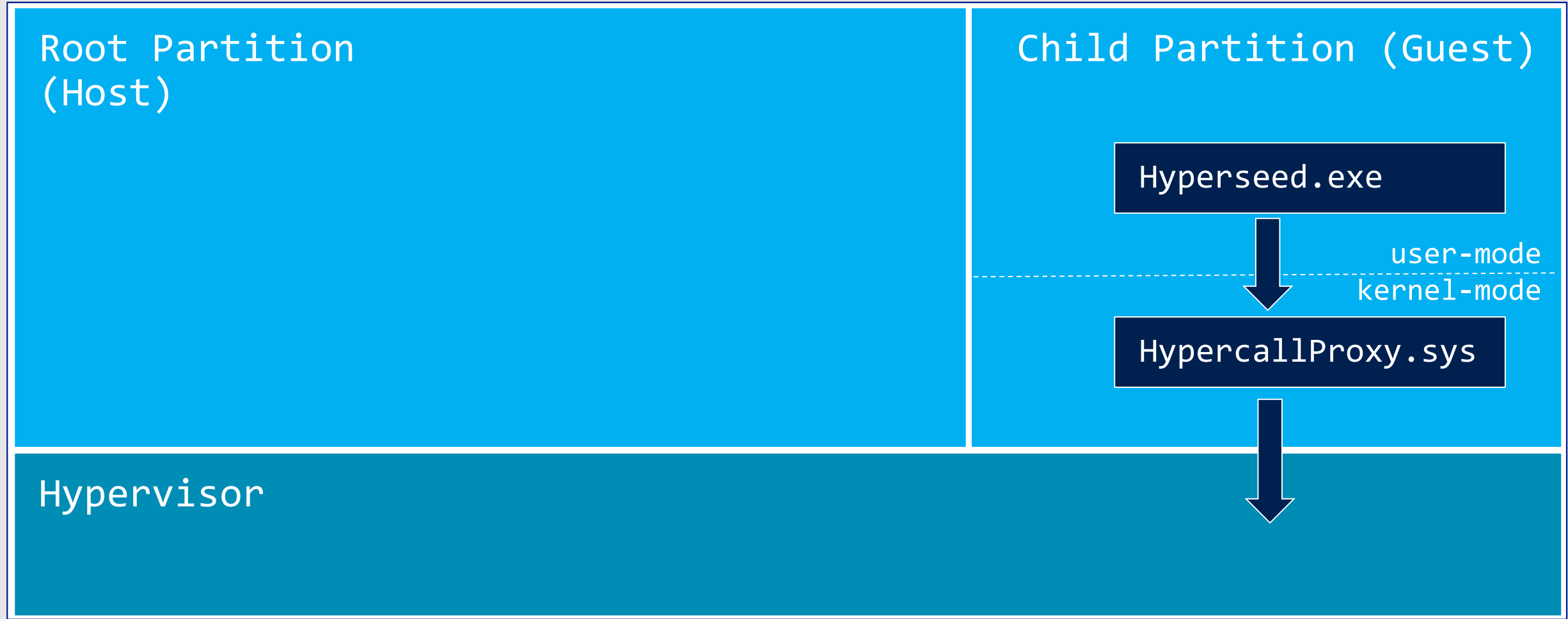
# Virtual Processor State Check

```
//  
// VP must be explicitly suspended before restore.  
//  
if (!Info->Vp->Activity.ExplicitSuspend)  
{  
    HT_ERR(status = HV_STATUS_INVALID_VP_STATE);  
}
```

# Hyperseed

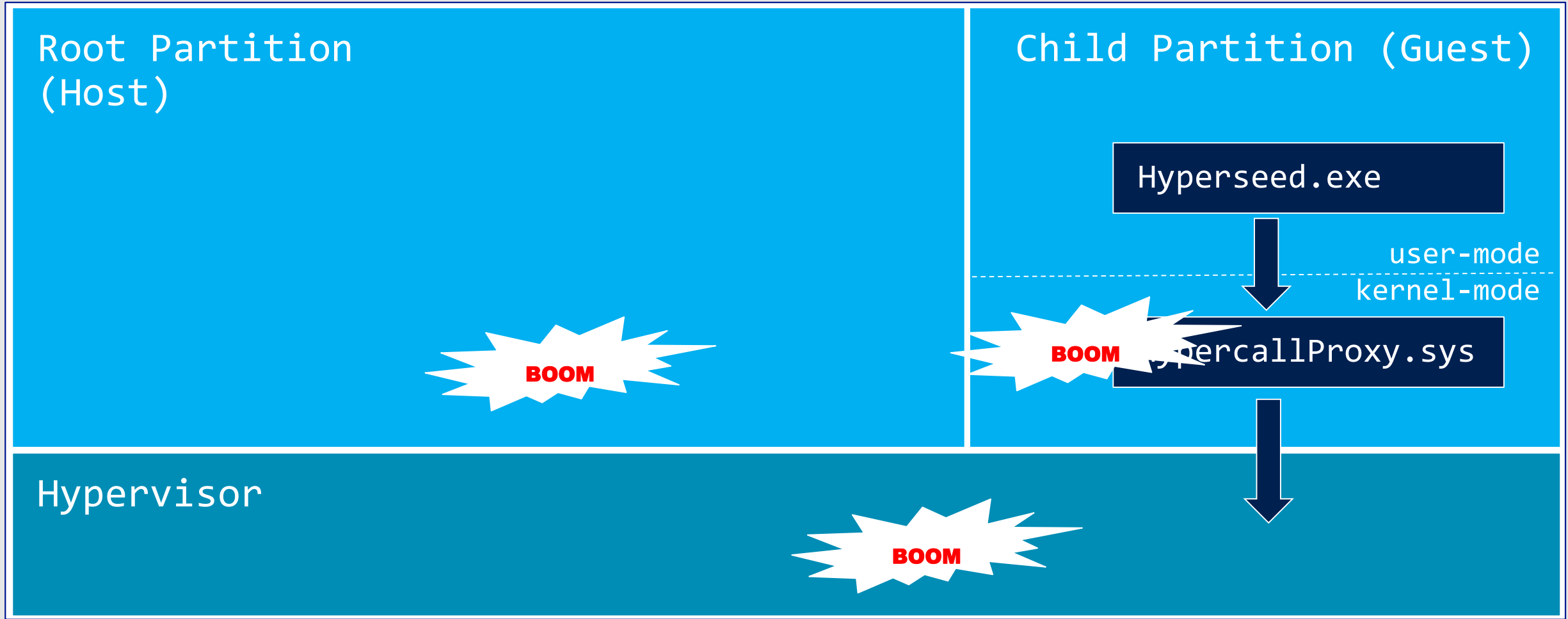


# Fuzzing from Guest

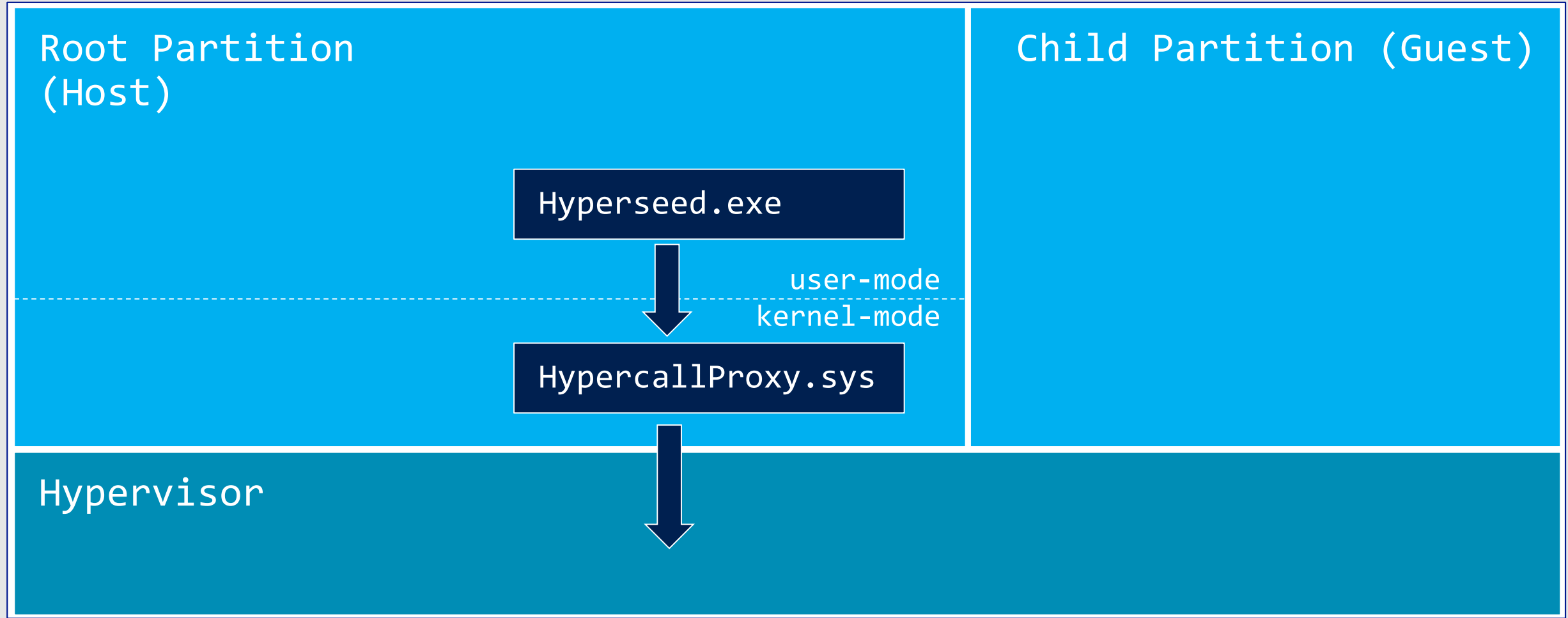




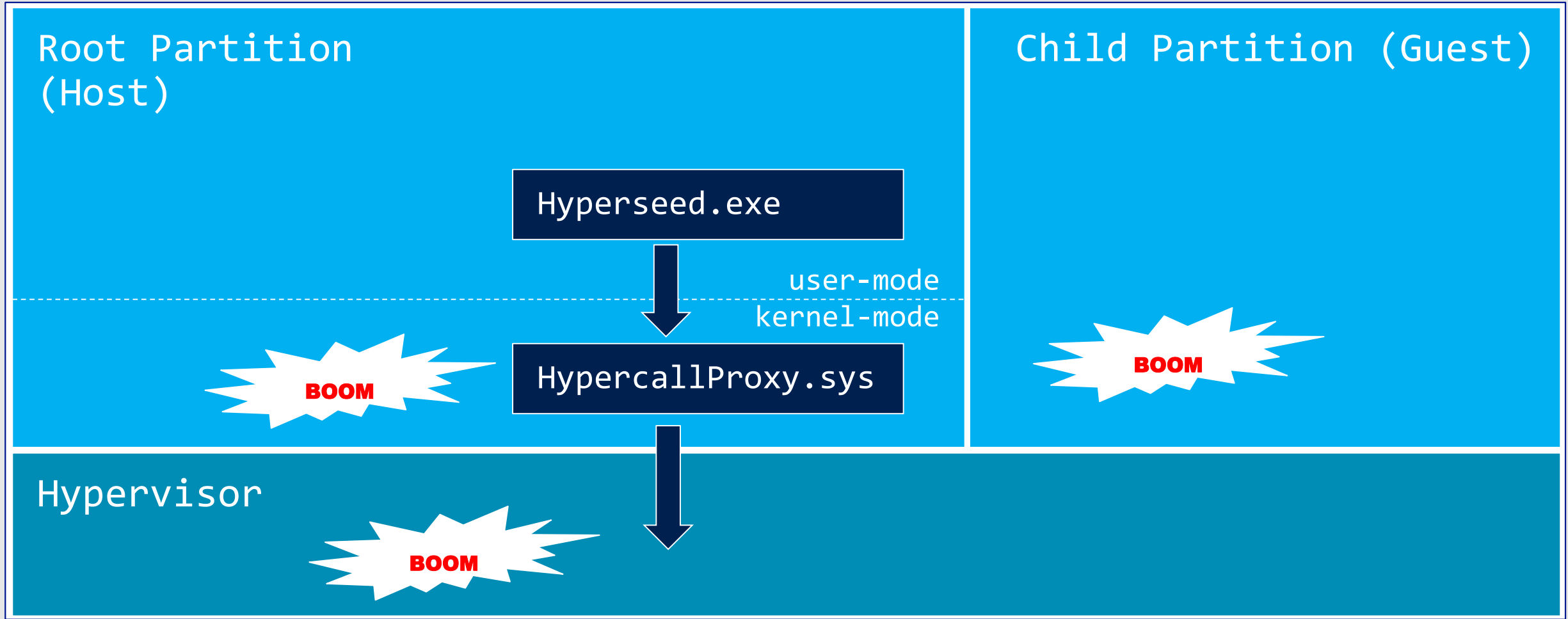
# Fuzzing from Guest



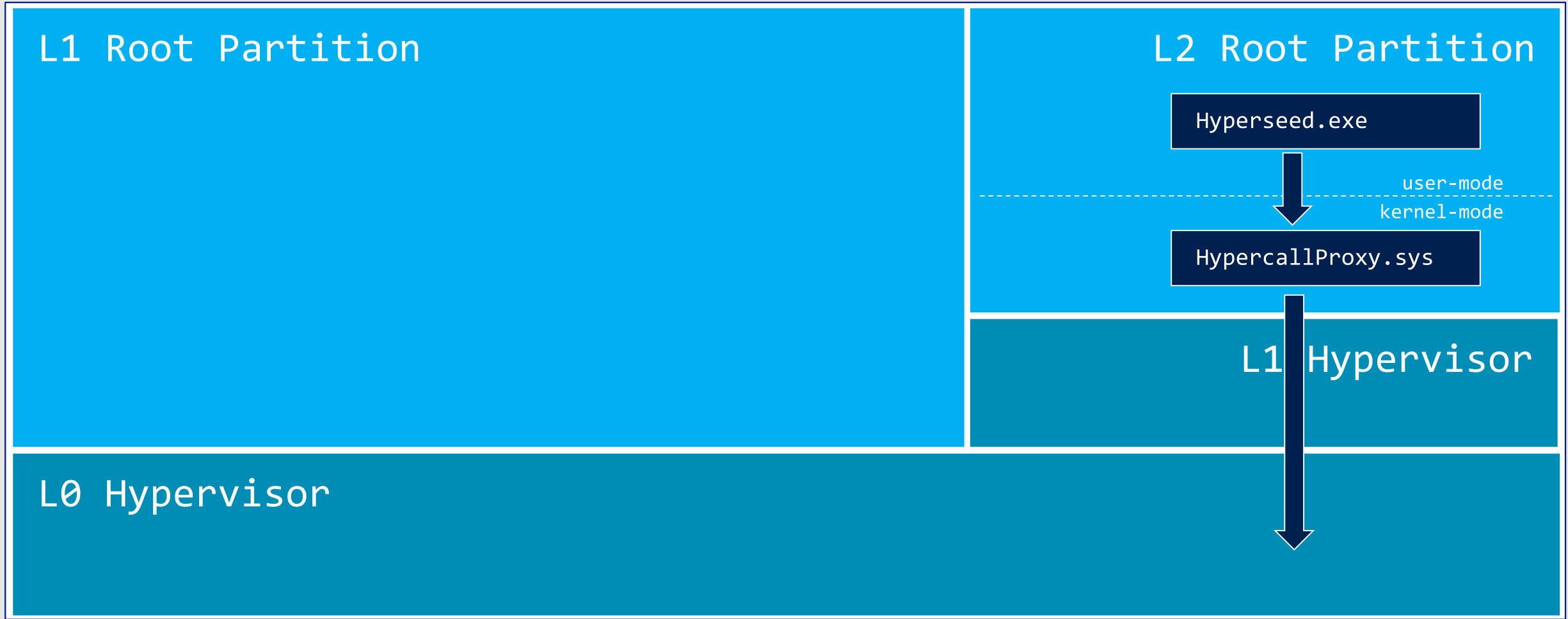
# Fuzzing from Root



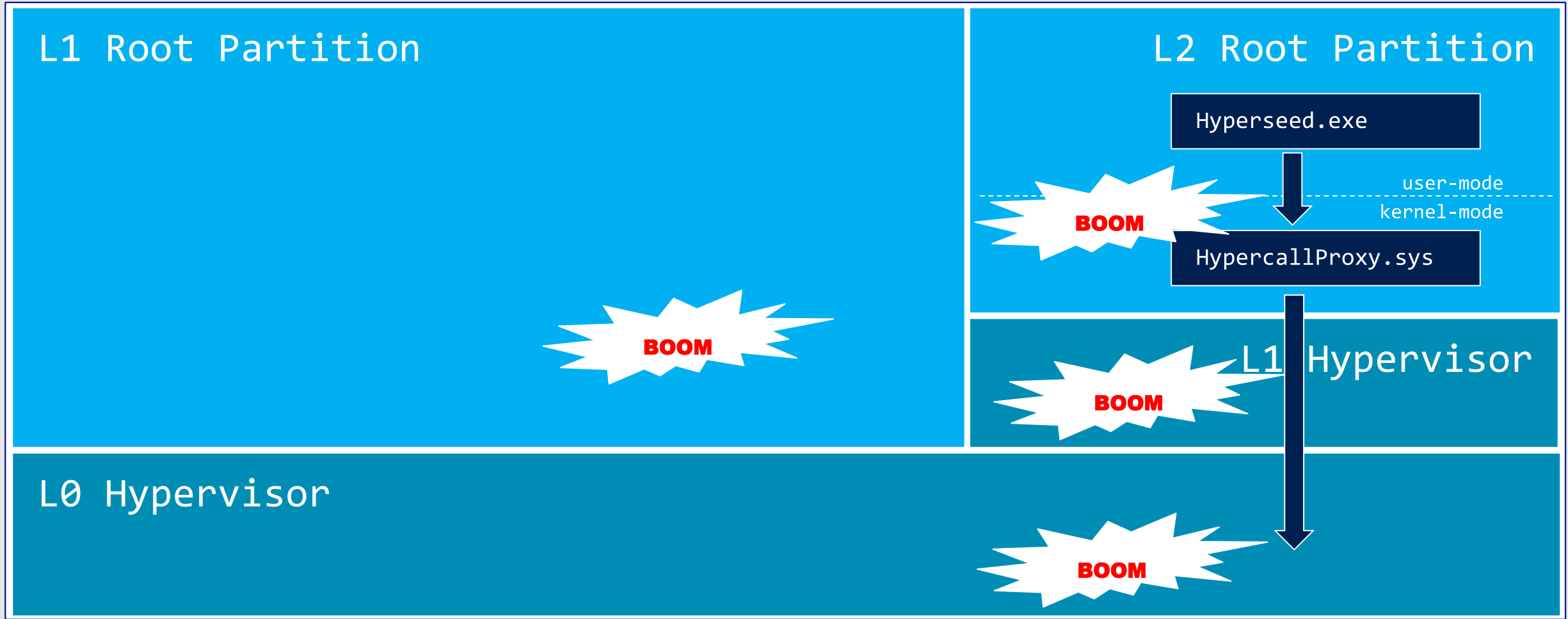
# Fuzzing from Root



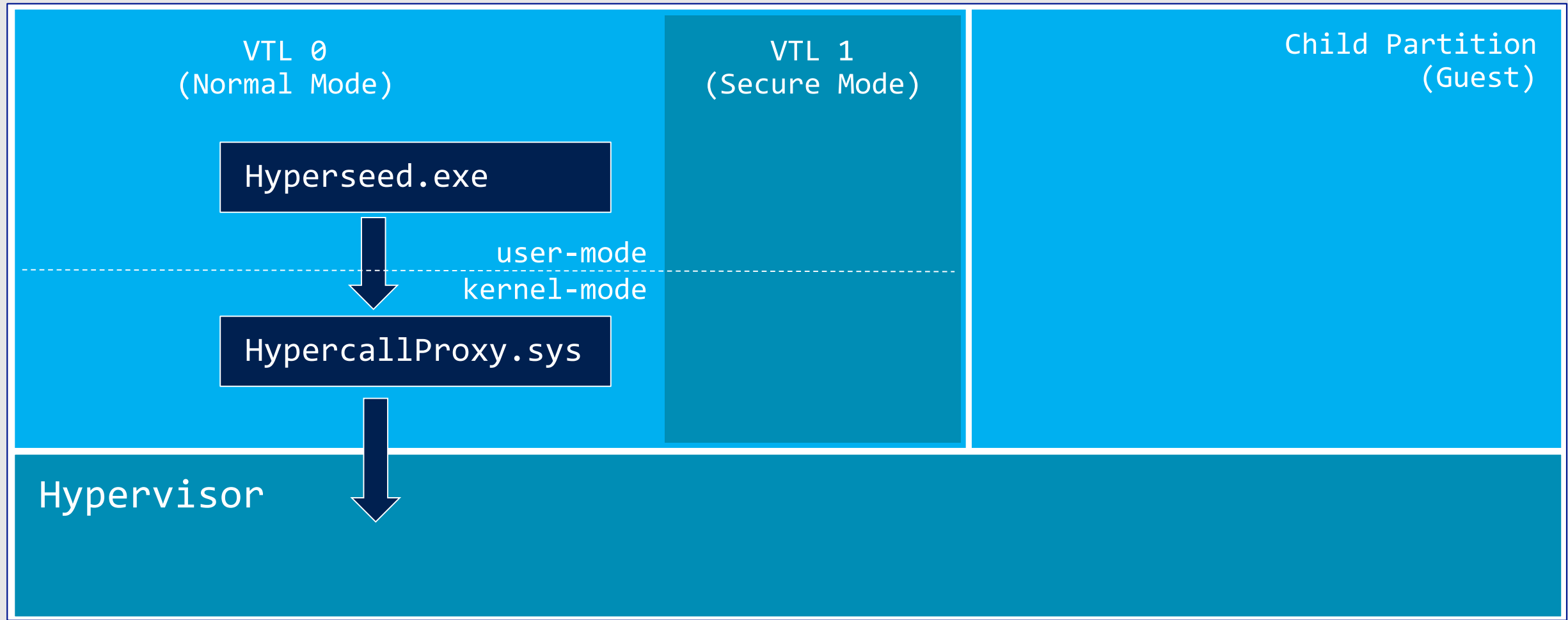
# Fuzzing from L2 Root Partition



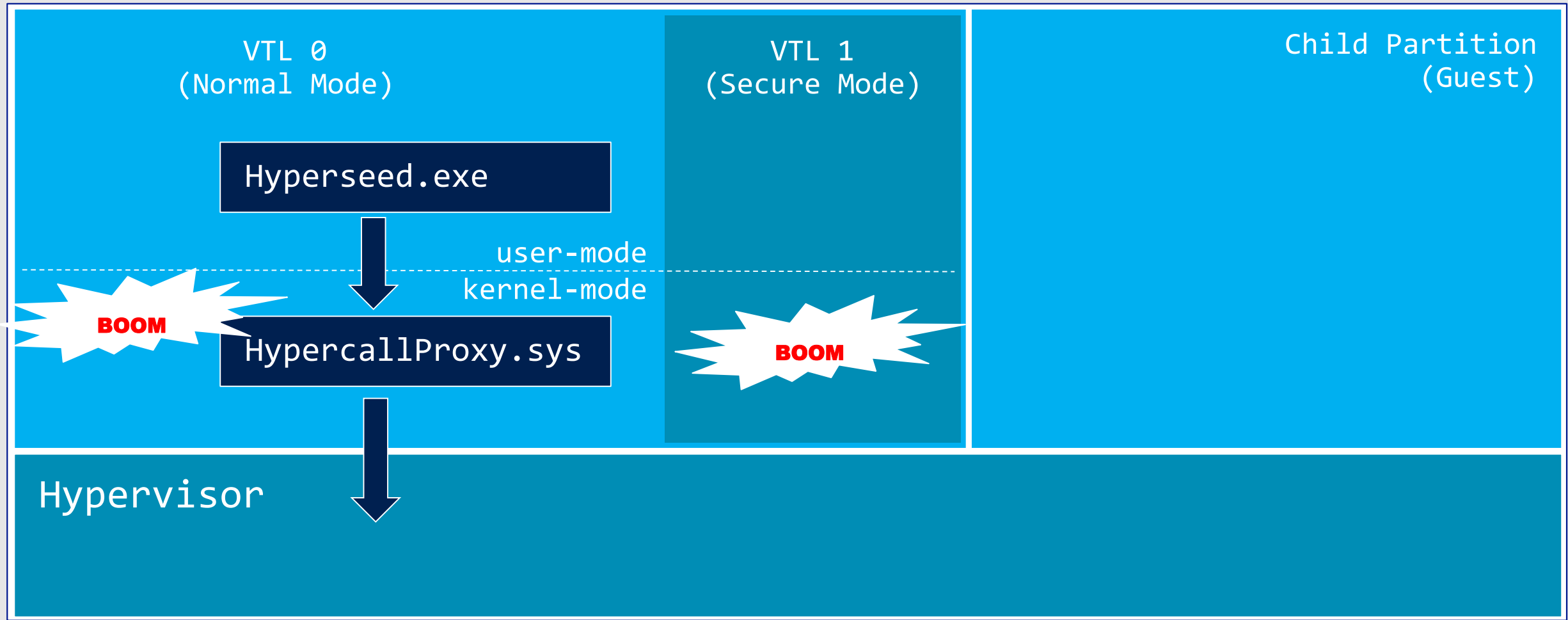
# Fuzzing from L2 Root Partition



# Fuzzing from Host VTL 0



# Fuzzing from Host VTL 0



# Format-Aware Fuzzing

## Past Fuzzing Experience

- CLFS, Deathnote of Microsoft Windows Kernel
- Intel PT & kAFL

## Pros

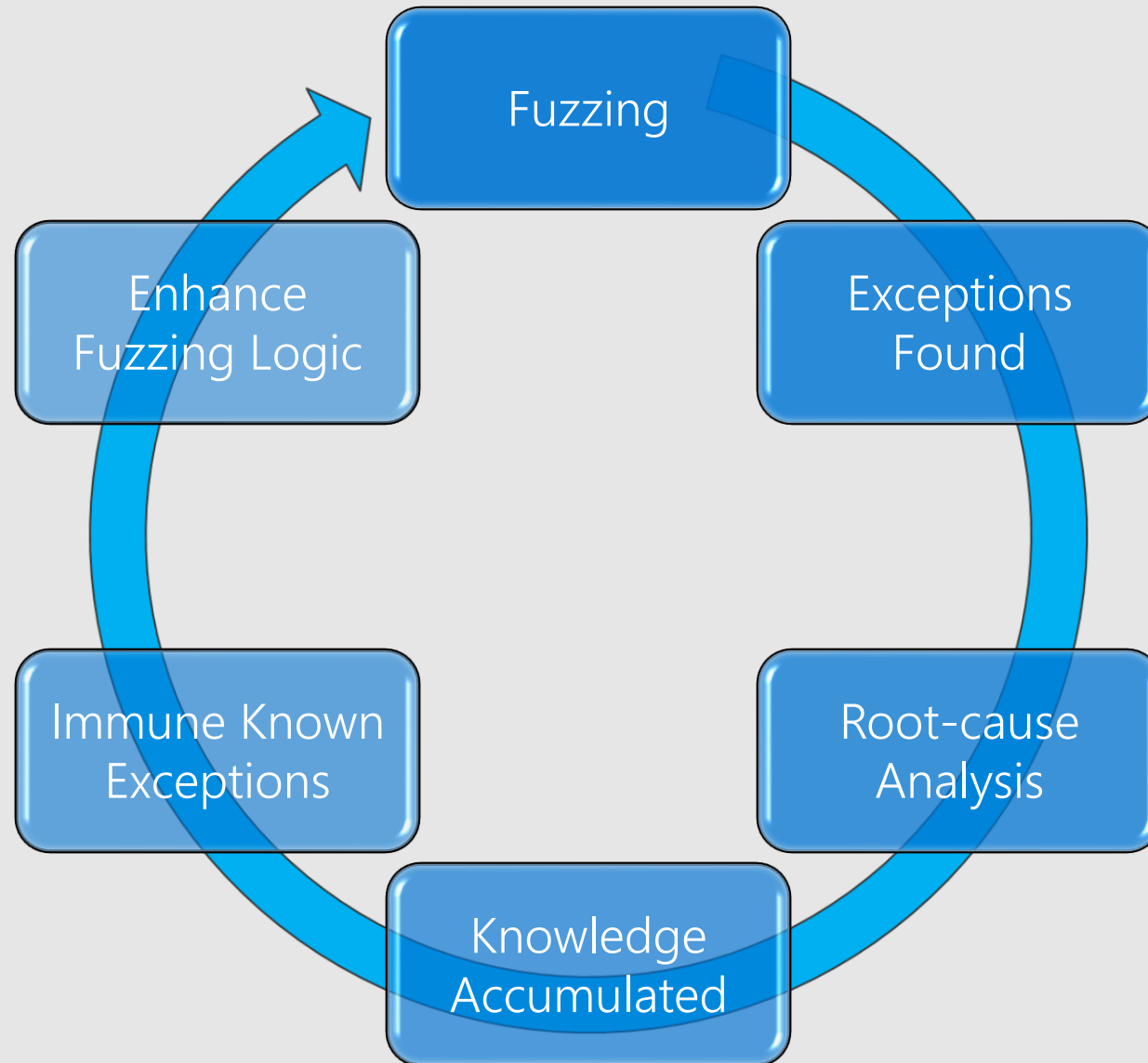
- Effective to bypass fields constraints
- Easy to scale, from coarse-grained to fine-grained, from less hypercalls to more hypercalls.

## Cons

- Assumptions may overlook some special circumstance
- Knowledge of input format needed



# Knowledge Accumulation Cycle



# How To Mutate?



# Primitive Data Types Mutator (uint64\_t)

**CRandomMutator:**

rand\_64() -> 0xFFFF123456789ABC

**CSlidingMutator:**

sliding(45) -> 0x13579FFFE2468ACF

**CSwapHalfMutator:**

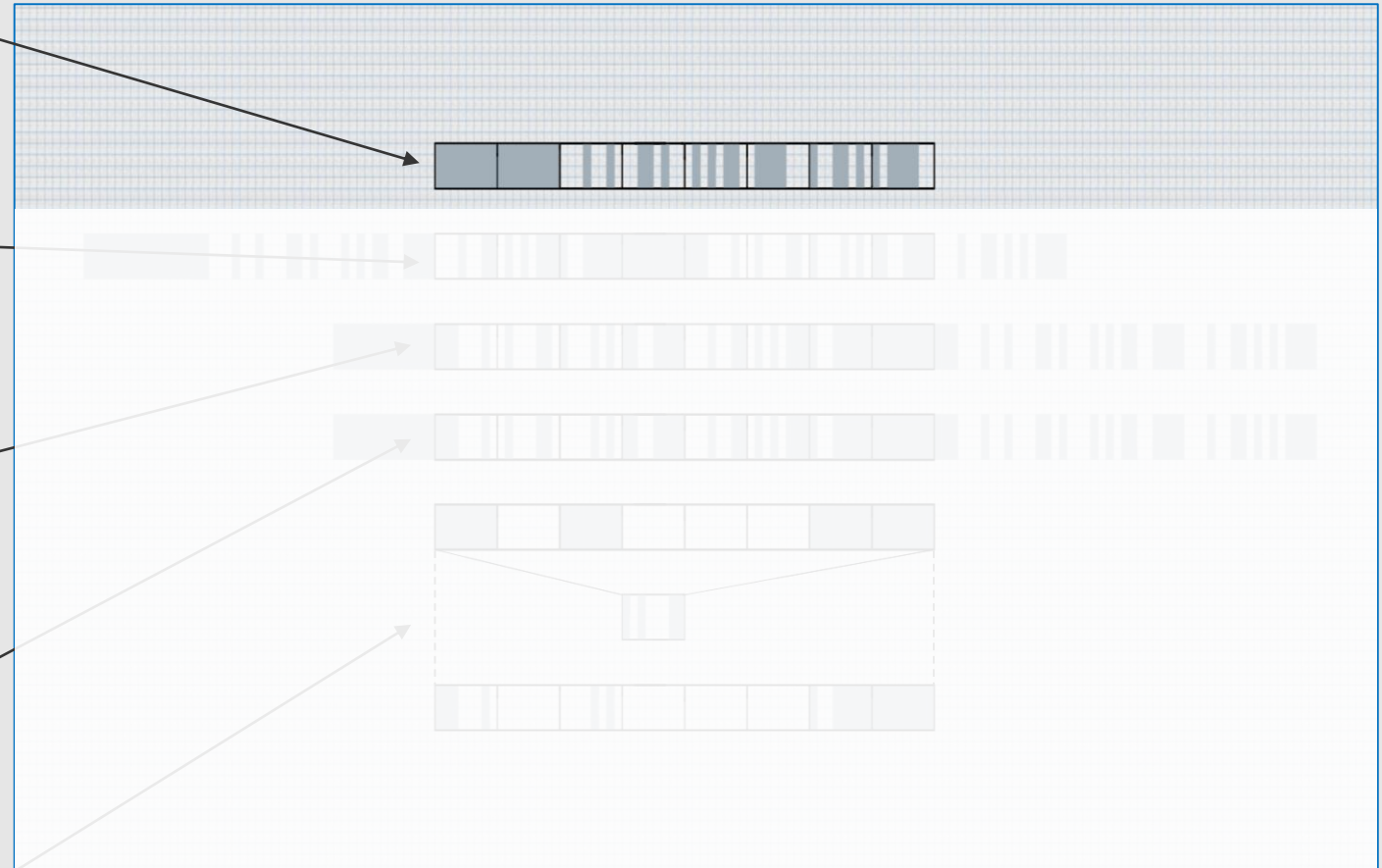
swap\_half() -> 0xE2468ACF13579FFF

**CBitFlipMutator:**

bitflip(47) -> 0xE2460ACF13579FFF

**CMaskMutator:**

mask(0xA3) -> 0xE2000A0000009FFF



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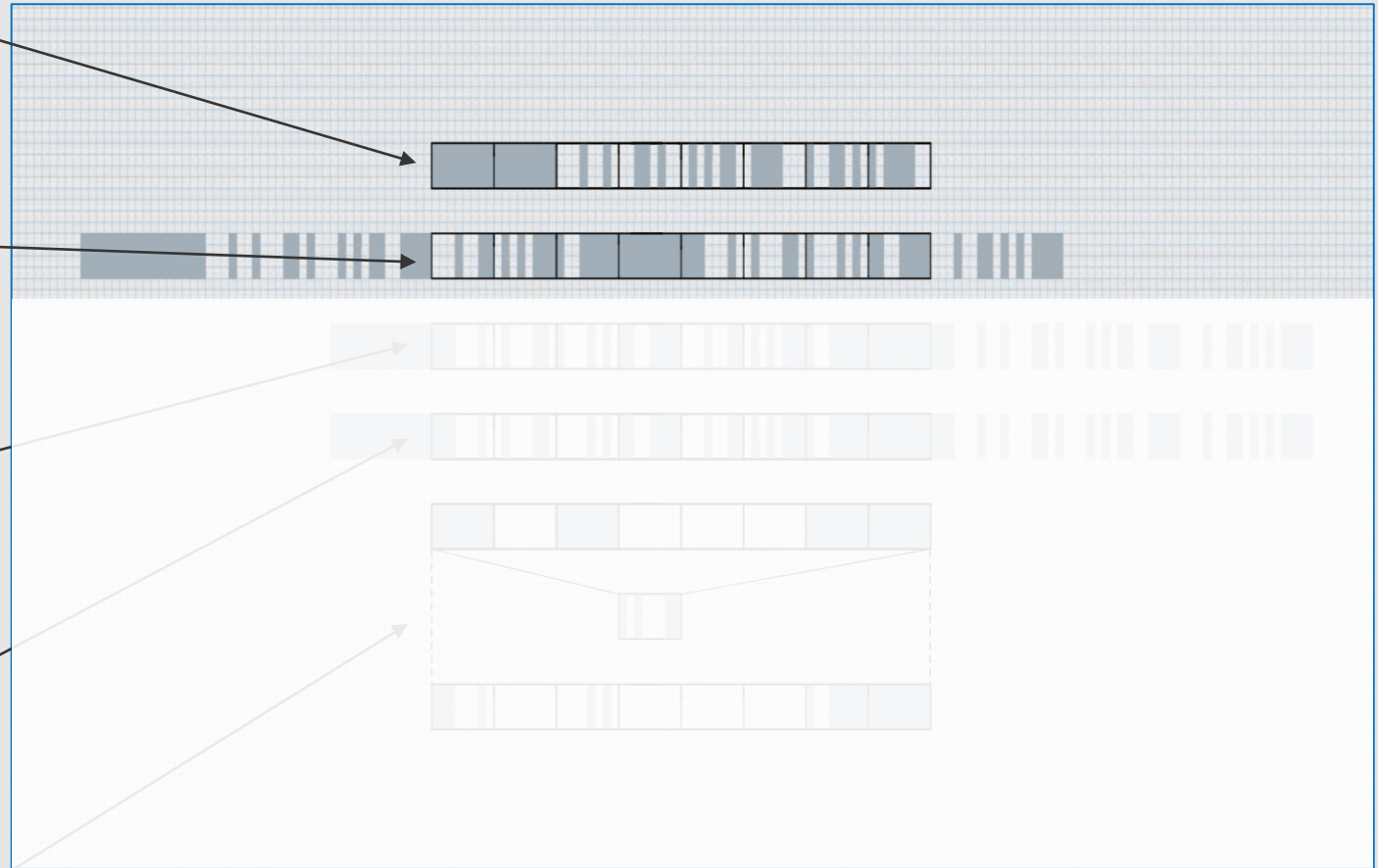
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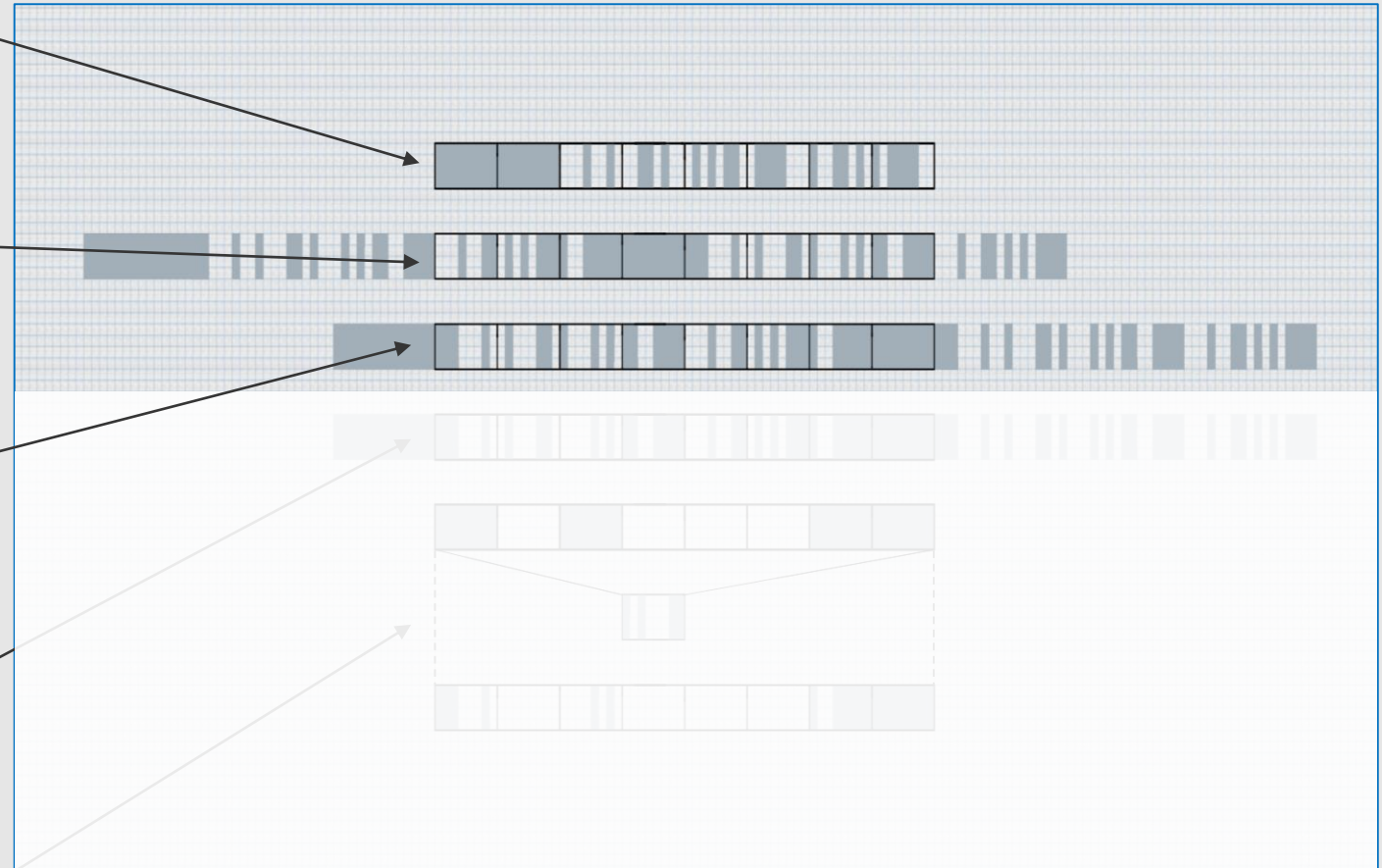
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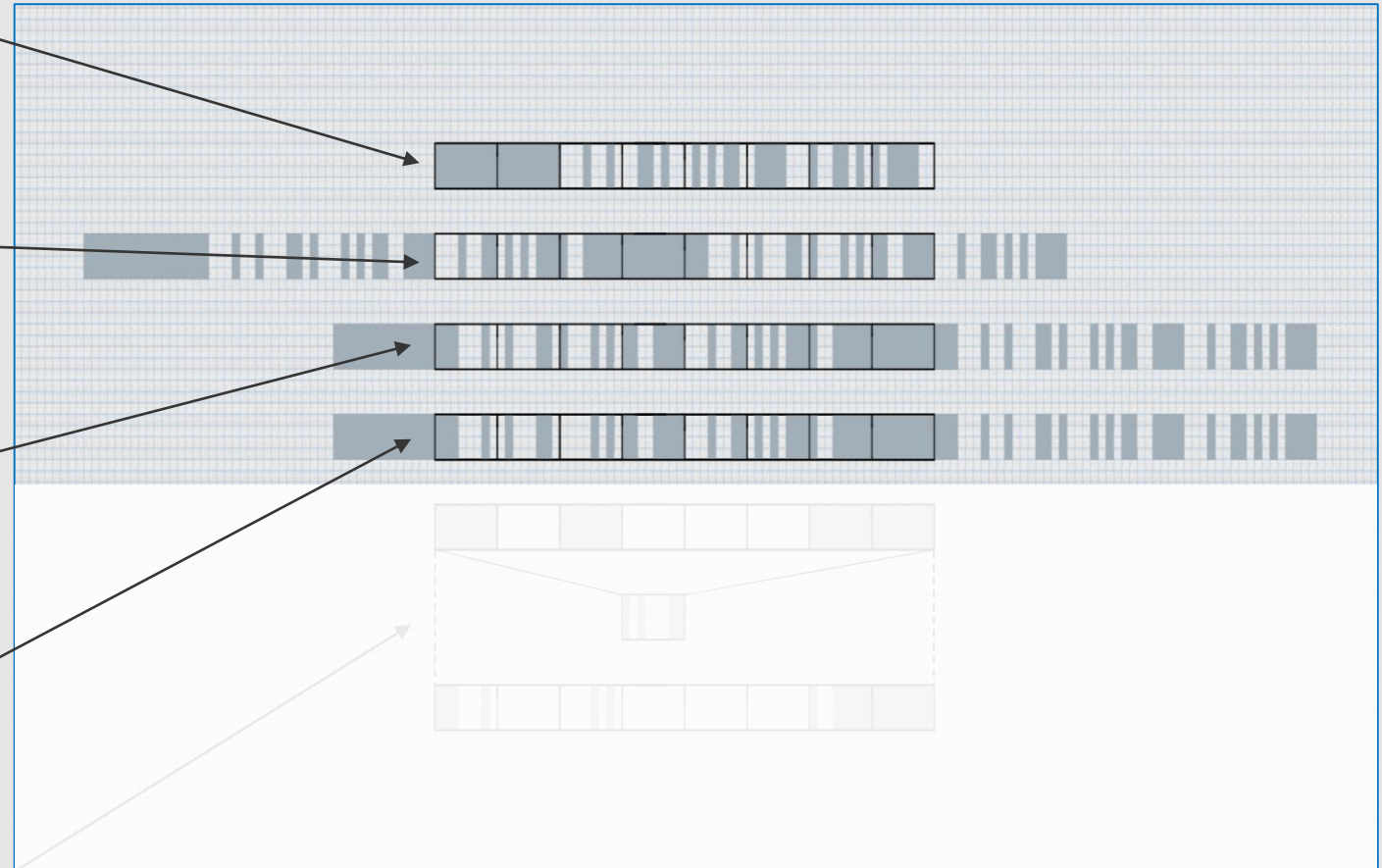
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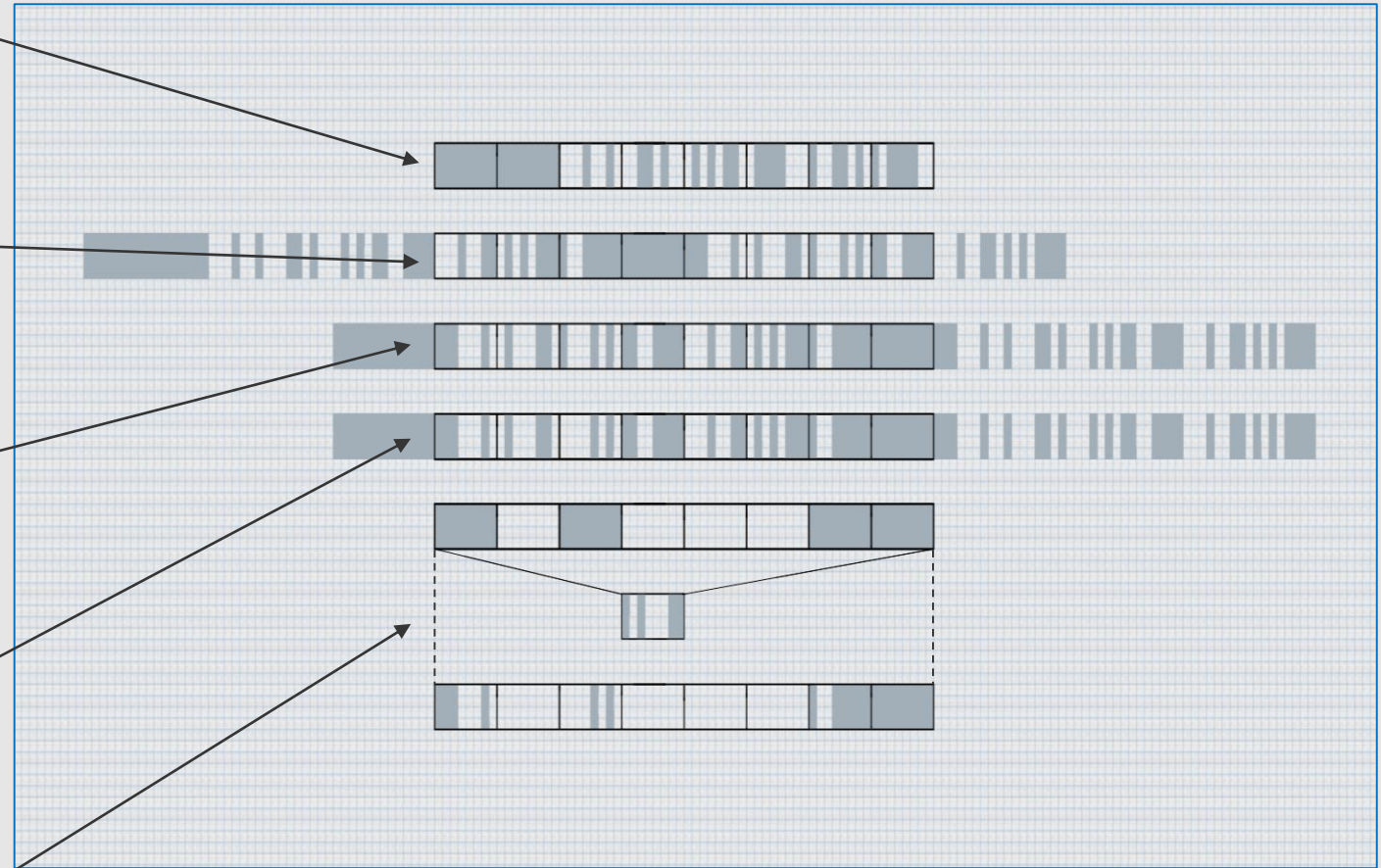
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# Mutator Dispatcher

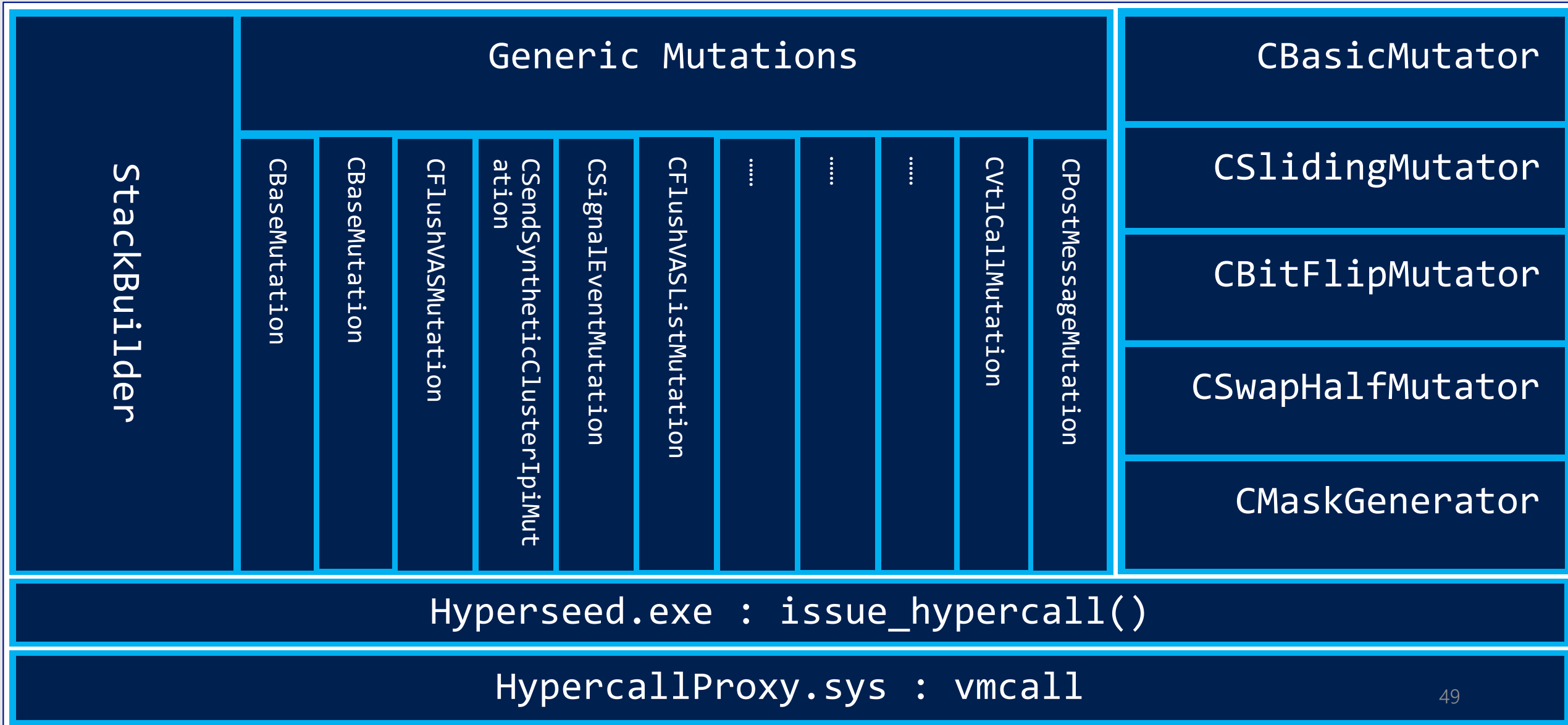
```
class CMutatorDispatcher
{
public:
    template<class T>
    static bool mutateBasicType(T& data, size_t hint = 0)
    {
        size_t bits = CPatternGenerator::Instance()->next_rand();
        size_t index = 0;

        if (is_bit_set(bits, index++)) CRandomMutator::mutateBasicType<T>(data, hint);
        if (is_bit_set(bits, index++)) CBitFlipMutator::mutateBasicType<T>(data, 1);
        if (is_bit_set(bits, index++)) CSlidingMutator::mutateBasicType<T>(data, hint);
        if (is_bit_set(bits, index++)) CBitFlipMutator::mutateBasicType<T>(data, 2);
        if (is_bit_set(bits, index++)) CSwapHalfMutator::mutateBasicType<T>(data, hint);
        if (is_bit_set(bits, index++)) CBitFlipMutator::mutateBasicType<T>(data, 4);
        if (is_bit_set(bits, index++)) CBitFlipMutator::mutateBasicType<T>(data, 8);
        if (is_bit_set(bits, index++)) CMaskMutator::mutateBasicType<T>(data, hint);

        return true;
    }
}
```



# Hyperseed Components



# CFlushVASMutation

## Wrapper Interface

```
HV_STATUS
HvFlushVirtualAddressSpace(
    __in HV_ADDRESS_SPACE_ID    AddressSpace,
    __in HV_FLUSH_FLAGS        Flags,
    __in UINT64 ProcessorMask
);
```

## Native Interface

### HvFlushVirtualAddressSpace

Call Code = 0x0002

### ➔ Input Parameters

0	AddressSpace (8 bytes)
8	Flags (8 bytes)
16	ProcessorMask (8 bytes)

```
#define SELECT          CPatternGenerator::Instance()->select
#define FROM_RANGE      CPatternGenerator::Instance()->from_range
#define RAND_8()        CPatternGenerator::Instance()->next_rand()
#define RAND_16()       CPatternGenerator::Instance()->next_rand_word()
#define RAND_32()       CPatternGenerator::Instance()->next_rand_dword()
#define RAND_64()       CPatternGenerator::Instance()->next_rand_qword()
#define BM(x)           CMutatorDispatcher::mutateBasicType(x)
#define BMN0(x)         CMutatorDispatcher::mutateBasicType_Nonzero(x)
#define DM(x, y)        CMutatorDispatcher::mutateData(x, y)
```

```
template<>
bool mutateCustom<HvCallFlushVirtualAddressSpaceTrait>(CBaseMutation* Mutation,
    HvCallFlushVirtualAddressSpaceBuffer& m_HypercallBuffer,
    shared_ptr<CBaseFeed> accum_feed)
{
    PHV_INPUT_FLUSH_VIRTUAL_ADDRESS_SPACE input = m_HypercallBuffer.InputBuffer();

    BM(input->Header.AddressSpace);
    // #define HV_FLUSH_ALL_PROCESSORS          (0x00000001)
    // #define HV_FLUSH_ALL_VIRTUAL_ADDRESS_SPACES (0x00000002)
    // #define HV_FLUSH_NON_GLOBAL_MAPPINGS_ONLY (0x00000004)
    // #define HV_FLUSH_USE_EXTENDED_RANGE_FORMAT (0x00000008)
    input->Header.Flags = RAND_8() & 0x07;
    BMN0(input->Header.ProcessorMask);

    accum_feed->set_input(string((char*)input, m_HypercallBuffer.input_len));
    return true;
}
```

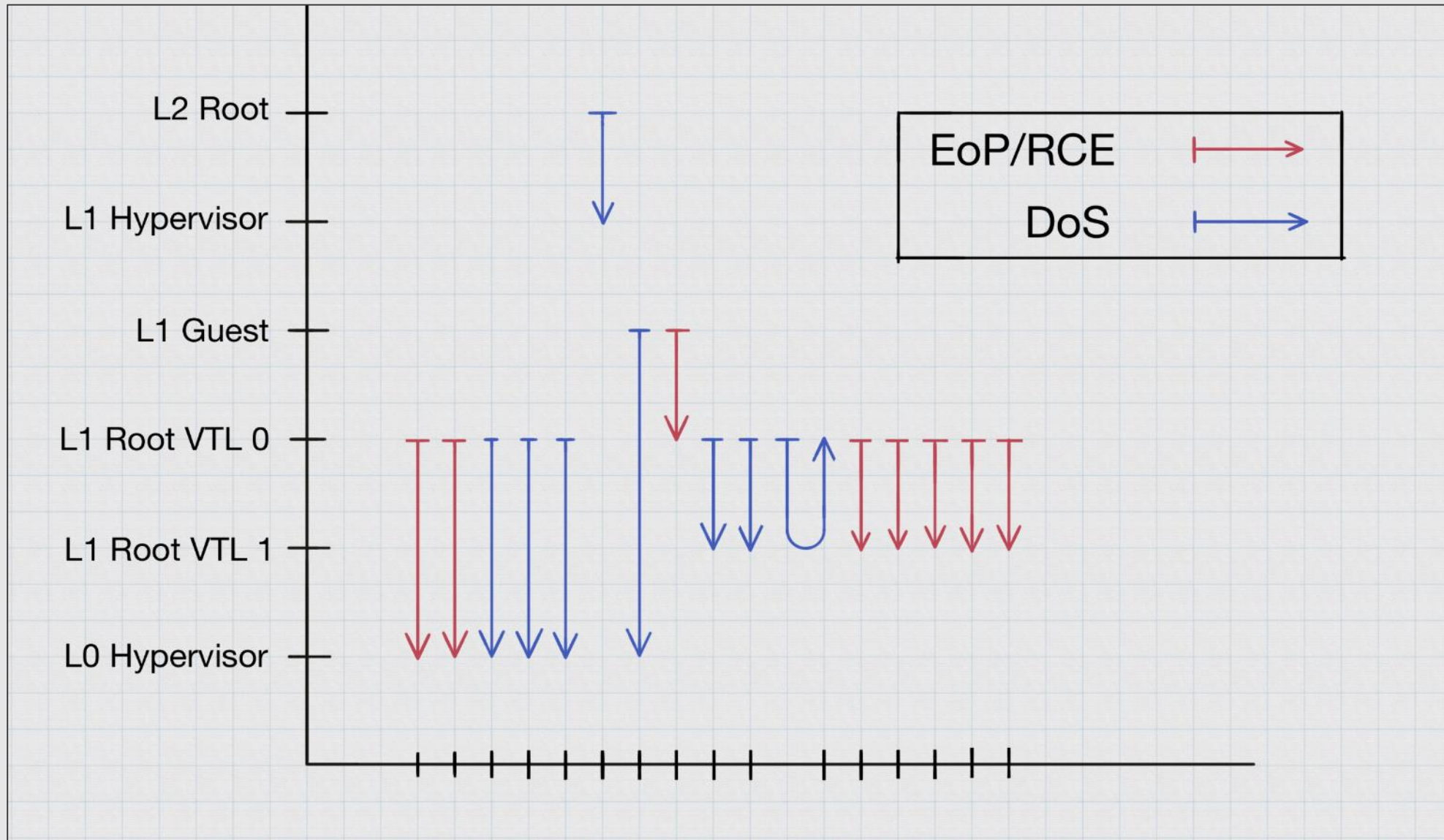
# Findings



# Fuzzing Stats

- Initiated 1 Year Ago
- 2 Contributors
- 350+ Source Files, 16K+ Lines of Code
- 169 Hypercalls Mutated
- 16 Findings So Far
  - 8 EoP/RCE
  - 8 DoS
  - 2 Qualified for Hyper-V Bounty
  - 7 Qualified for WIP Bounty

# Findings

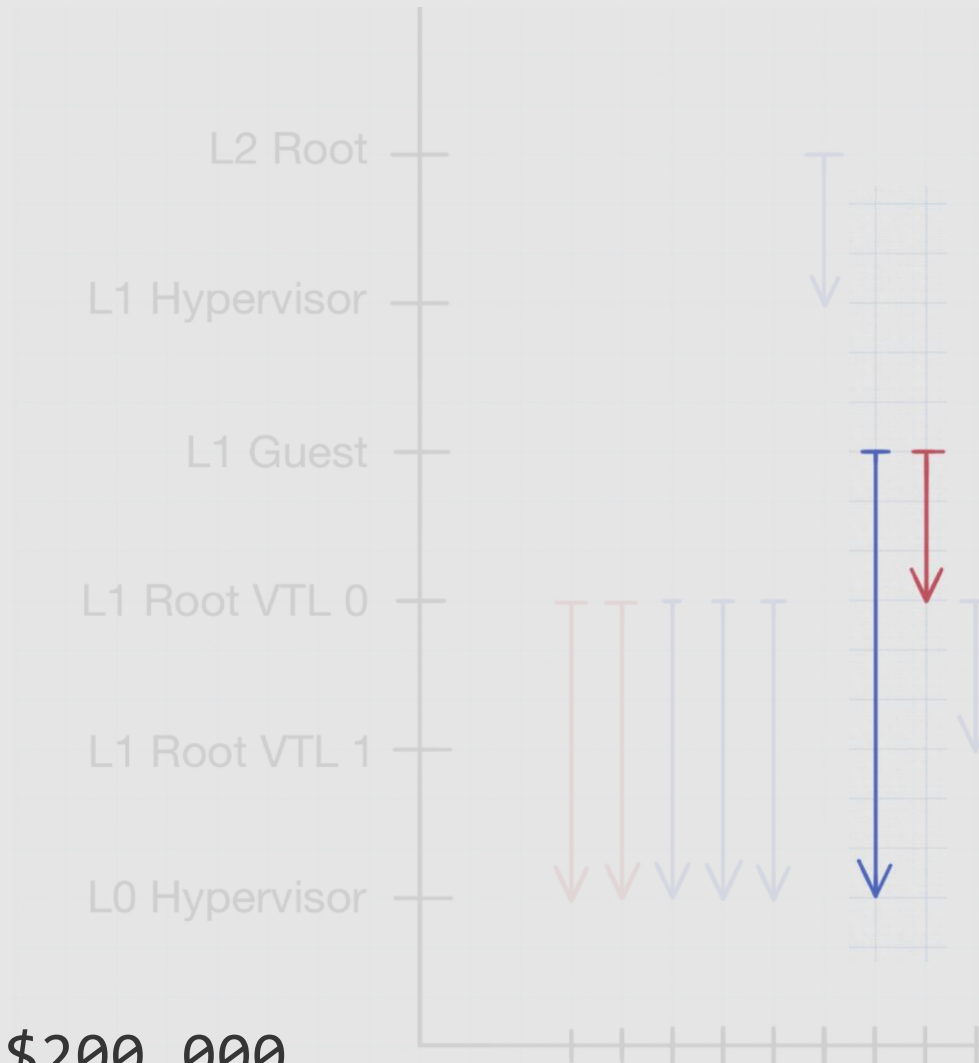


# How Much Are They Worth? (Hyper-V Bounty)



✓ 1 x \$200,000

✓ 1 x \$15,000



FROM: <https://www.microsoft.com/en-us/msrc/bounty-hyper-v>

## Remote Code Execution

An eligible submission includes a RCE vulnerability in Microsoft Hyper-V that enables a guest virtual machine to compromise the hypervisor, escape from a guest virtual machine to the host, or escape from one guest virtual machine to another guest virtual machine.

Vulnerability Type	Tier	Proof of concept	Functioning Exploit	Report Quality	Payout range (USD)*
RCE	Tier 1	Required	Yes	High	\$250,000
			No	High	\$200,000
			No	Low	\$50,000
RCE	Tier 2	Required	Yes	High	\$150,000
			No	High	\$100,000
			No	Low	\$25,000
RCE	Tier 3	Required	Yes	High	\$20,000
			No	High	\$15,000
			No	Low	\$5,000

## Denial of Service and Information Disclosure

The vulnerability should result in one of the following:

- Crash the host machine, resulting in a denial of service condition
- Cause a failure to start and stop VMs
- Gain sensitive information from the host machine or another guest

Vulnerability Type	Tier	Proof of concept	Report Quality	Payout range (USD)
DOS	Tier 1	Required	High	\$15,000
			Low	\$5,000
Info Disclosure	Tier 1	Required	High	\$25,000
			Low	\$5,000
	Tier 2	Required	High	\$15,000
			Low	\$5,000

# How Much Are They Worth? (WIP Bounty)

FROM: <https://www.microsoft.com/en-us/msrc/bounty-windows-insider-preview>

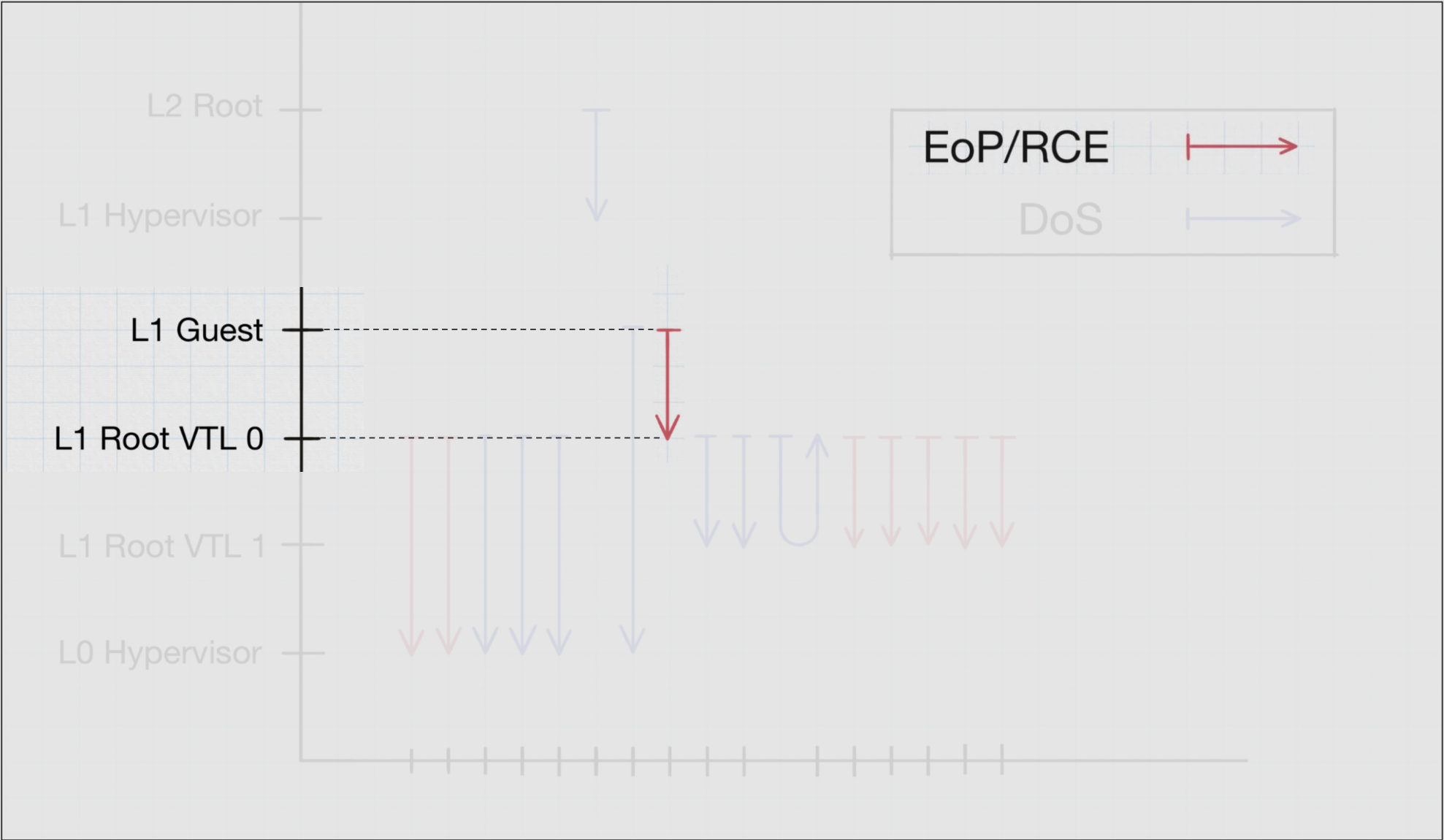
Impact ↑↓	Report Quality	Severity			
		Critical	Important	Moderate	Low
Remote Code Execution	High	\$30,000	\$25,000	N/A	N/A
	Medium	\$12,000	\$8,000		
	Low	\$3,000	\$2,000		
Elevation of Privilege	High	N/A	\$20,000	\$0	N/A
	Medium		\$8,000		
	Low		\$2,000		
Security Feature Bypass	High	N/A	\$20,000	\$0	\$0
	Medium		\$8,000		
	Low		\$2,000		
Information Disclosure	High	N/A	\$10,000	\$0	\$0
	Medium		\$3,000		
	Low		\$1,000		



✓ 2 x \$20,000

✓ 5 x \$20,000

# CVE-2018-8439





# CVE-2018-8439

FROM: <https://portal.msrc.microsoft.com/en-US/security-guidance/advisory/CVE-2018-8439>

## CVE-2018-8439 | Windows Hyper-V Remote Code Execution Vulnerability

### Security Vulnerability

Published: 09/11/2018  
[MITRE CVE-2018-8439](#)

A remote code execution vulnerability exists when Windows Hyper-V on a host server fails to properly validate input from an authenticated user on a guest operating system. To exploit the vulnerability, an attacker could run a specially crafted application on a guest operating system that could cause the Hyper-V host operating system to execute arbitrary code.

An attacker who successfully exploited the vulnerability could execute arbitrary code on the host operating system.

The security update addresses the vulnerability by correcting how Hyper-V validates guest operating system user input.

### Exploitability Assessment

The following table provides an [exploitability assessment](#) for this vulnerability at the time of original publication.

Publicly Disclosed	Exploited	Latest Software Release	Older Software Release	Denial of Service
No	No	2 - Exploitation Less Likely	2 - Exploitation Less Likely	Not Applicable

### On this page

- [Executive Summary](#)
- [Exploitability Assessment](#)
- [Security Updates](#)
- [Mitigations](#)
- [Workarounds](#)
- [FAQ](#)

### Acknowledgements

Daniel King ([@long123king](#)), MSRC<sup>57</sup> Microsoft

# CVE-2018-8439

## HvPostMessage

- The hypervisor provides a simple inter-partition communication facility that allows one partition to send a parameterized message to another partition. (Because the message is sent asynchronously, it is said to be posted.) The destination partition may be notified of the arrival of this message through an SynIC(Synthetic Interrupt Controller) interrupt.

FROM: <https://docs.microsoft.com/en-us/virtualization/hyper-v-on-windows/reference/tlfs>

Wrapper Interface

HV\_STATUS

HvPostMessage(

\_\_in HV\_CONNECTION\_ID    ConnectionId,

\_\_in HV\_MESSAGE\_TYPE    MessageType,

\_\_in UINT32 PayloadSize,

\_\_in\_ecount(PayloadSize)

PCVOID Message

);

Native Interface

HvPostMessage

Call Code = 0x005C

➔ Input Parameters

0	ConnectionId (4 bytes)	Padding (4 bytes)
8	MessageType (4 bytes)	PayloadSize (4 bytes)
16	Message[0] (8 bytes)	
⋮	⋮	
248	Message[29] (8 bytes)	

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# CVE-2018-8439

- Reproduce Steps

- I. Run Hyperseed inside guest VM to keep fuzzing HvPostMessage with ChannelMessageOpenChannel(0x05)/ChannelMessageCloseChannel(0x07) continuously.
- II. Manually reset guest VM from host.

- Result

- I. Host BSOD

# CVE-2018-8439

```
0: kd> !analyze -v
*****
*                                     *
*                               Bugcheck Analysis                               *
*                                     *
*****

DRIVER_IRQL_NOT_LESS_OR_EQUAL (d1)
An attempt was made to access a pageable (or completely invalid) address at an
interrupt request level (IRQL) that is too high.  This is usually
caused by drivers using improper addresses.
If kernel debugger is available get stack backtrace.
Arguments:
Arg1: fffffba81d8001008, memory referenced
Arg2: 0000000000000002, IRQL
Arg3: 0000000000000001, value 0 = read operation, 1 = write operation
Arg4: fffff80fb4b51122, address which referenced memory
0: kd> .trap 0xfffff58dbc6ca7a0
NOTE: The trap frame does not contain all registers.
Some register values may be zeroed or incorrect.
rax=fffff8780eb36adc0 rbx=0000000000000000 rcx=0000000000000000
rdx=fffffba81d8001008 rsi=0000000000000000 rdi=0000000000000000
rip=fffff80fb4b51122 rsp=fffff58dbc6ca930 rbp=0000000000000005
r8=0000000000000001 r9=0000000000000000 r10=000000000000000a
r11=fffffba81d7a6a010 r12=0000000000000000 r13=0000000000000000
r14=0000000000000000 r15=0000000000000000
iopl=0         nv up ei ng nz na pe nc
vmbusr!BusChSendInterrupt+0x42:
fffff80f`b4b51122 f00fab0a  lock bts dword ptr [rdx],ecx ds:fffffba81`d8001008=????????
0: kd> kf
*** Stack trace for last set context - .thread/.cxr resets it
#  Memory  Child-SP      RetAddr      Call Site
00  fffff58d`bc6ca930 fffff80f`b5542e5c vmbusr!BusChSendInterrupt+0x42
01  30 fffff58d`bc6ca960 fffff80f`b55429a5 vmbkmclr!InCompletePacket+0x4ac
02  c0 fffff58d`bc6caa20 fffff80f`b72f112b vmbkmclr!VmbChannelPacketComplete+0x15
03  30 fffff58d`bc6caa50 fffff80f`b55413a2 storvsp+0x112b
04  60 fffff58d`bc6caab0 fffff80f`b4b55680 vmbkmclr!InpProcessingWorkerRoutine+0x212
05  80 fffff58d`bc6cab30 fffff800`d28fdae7 vmbusr!AwWorkerThread+0xb0
06  60 fffff58d`bc6cab90 fffff800`d29bbb86 nt!PspSystemThreadStartup+0x47
07  50 fffff58d`bc6cabe0 00000000`00000000 nt!KiStartSystemThread+0x16
0: kd> !pte rdx
```

```
VA fffffba81d8001008
PXE at FFFF98CC66331BA8  PPE at FFFF98CC66375038  PDE at FFFF98CC6EA07600
contains 0A0000000515D863 contains 0A0000000515E863 contains 0A00000023ACB5863
pfn 515d  ---DA--KWEV  pfn 515e  ---DA--KWEV  pfn 23acb5  ---DA--KWEV
```

```
0: kd> dt vmbusr!_HV_MONITOR_PAGE -b
+0x000 TriggerState      : _HV_MONITOR_TRIGGER_STATE
+0x000 AsUINT32          : Uint4B
+0x000 GroupEnable       : Pos 0, 4 Bits
+0x000 RsvdZ             : Pos 4, 28 Bits
+0x004 RsvdZ1            : Uint4B
+0x008 TriggerGroup      : _HV_MONITOR_TRIGGER_GROUP
+0x000 AsUINT64          : Uint8B
+0x000 Pending           : Uint4B
+0x004 Armed             : Uint4B
+0x028 RsvdZ2            : Uint8B
+0x040 NextCheckTime     : Int4B
+0x240 Latency           : Uint2B
+0x340 RsvdZ3            : Uint8B
+0x440 Parameter         : _HV_MONITOR_PARAMETER
+0x000 ConnectionId      : _HV_CONNECTION_ID
+0x000 AsUINT32          : Uint4B
+0x000 Id                : Pos 0, 24 Bits
+0x000 Reserved          : Pos 24, 8 Bits
+0x004 FlagNumber        : Uint2B
+0x006 RsvdZ             : Uint2B
+0x840 RsvdZ4            : UChar
```

```
PTE at FFFF98DD40EC0008
contains 001D845D00000000
not valid
Page has been freed
```

# CVE-2018-8439

- Monitored Notification Page
  - Allocated in `vmbusr!ParentCreateMonitors`
    - When initializing VMBus connection from the client
  - Freed in `vmbusr!ParentDeleteMonitors`
    - When resetting, destroying partition or failing VMBus connection
  - Cache a reference in `vmbusr!PncEnableInterrupt`
    - `BusFdoOpenChannel`, holds the reference even if the channel is not opened successfully
  - Clear the cached reference in `vmbusr!PncDisableInterrupt`
    - `BusFdoCloseChannel`, won't call this routine if it is not opened successfully
  - Accessing the cached reference after being freed leads to UAF

# CVE-2018-8439 (Fix)

- Monitored Notification Page
  - Allocated in `vmbusr!ParentCreateMonitors`
    - When initializing VMBus connection
  - Freed in `vmbusr!ParentDeleteMonitors`
    - When resetting, destroying partition or failing VMBus connection
- Cache a reference in `vmbusr!PncEnableInterrupt`
  - `BusFdoOpenChannel`, holds the reference even if the channel not opened successfully
- Clear the cached reference in `vmbusr!PncDisableInterrupt`
  - `BusFdoCloseChannel`, won't be called if the channel is not opened successfully
  - `BusFdoOpenResult`, clear the cached reference if the channel is not opened properly.

# Future



# Future

- Keep Improving Format-Aware Fuzzing
- Adopt Coverage-Guided Fuzzing
  - Intel PT & AFL
- More Details on Findings After Being Patched
- Open-Sourcing
- More Targets



# Acknowledgements

- Microsoft Security Response Center
- CDG Platform Security & Vulnerability Research
- Hyper-V Development Team

# Appendix

# Past Research

- [Security Assessment of Microsoft Hyper-V \(ERNW Newsletter 43 / May 2014, Felix Wilhelm, Mattias Luft\)](#)
- [Ring 0 to Ring -1 \(Syscan 2015 – Alex Ionescu\)](#)
- [The Battle of SKM and IUM \(Blackhat 2015 - Alex Ionescu\)](#)
- [Analysis of the Attack Surface of Windows 10 Virtualization-Base Security \(Blackhat 2016 - Rafal Wojtczuk\)](#)
- [VBS and VSM Internals \(Saar Amar\)](#)
- [A Dive in to Hyper-V Architecture and Vulnerabilities \(Blackhat 2018 - Joe Bialek & Nicolas Joly\)](#)
- [Hardening Hyper-V through offensive security research \(Blackhat 2018 – Jordan Rabet\)](#)
- [First steps in Hyper-V research \(Saar Amar\)](#)
- [Fuzzing para-virtualized devices in Hyper-V \(MSFT Virtualization Security Team\)](#)
- [Writing a Hyper-V “Bridge” for Fuzzing \(Alex Ionescu\)](#)

# Useful Public Resources

- Hypervisor Top-Level Functional Specification
  - <https://docs.microsoft.com/en-us/virtualization/hyper-v-on-windows/reference/tlfs>
- Hyper-V Debugging Symbols
  - <https://blogs.technet.microsoft.com/virtualization/2018/04/25/hyper-v-symbols-for-debugging/>
- Linux Integration Service for Microsoft Hyper-V
  - <https://github.com/LIS>
- Unofficial hdk – Hyper-V Development Kit
  - <https://ionescu007.github.io/hdk/>
- Report quality definitions for Microsoft's Bug Bounty programs
  - <https://www.microsoft.com/en-us/msrc/bounty-example-report-submission>

Q&A



